



University of
Massachusetts
Amherst

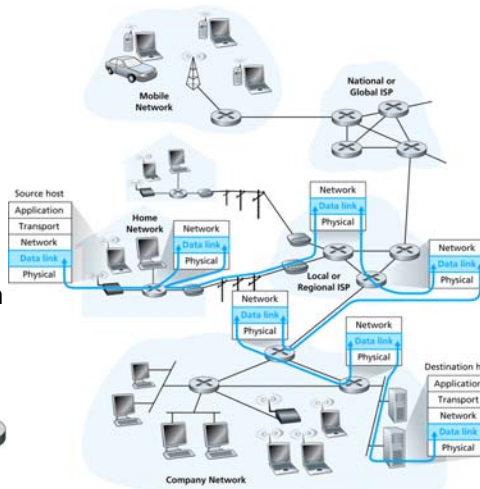
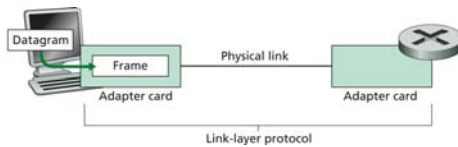
ECE697AA – Lecture 5

Link Layer: Ethernet

Tilman Wolf
Department of Electrical and Computer Engineering
09/16/08

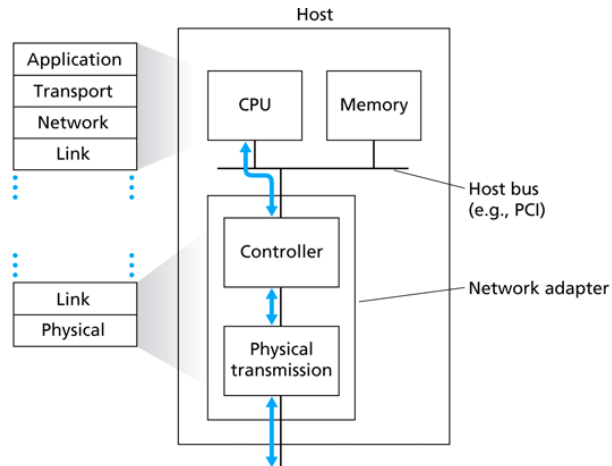
Data link layer

- Data links connect one device to another
 - Single hop connection
- Services
 - Framing
 - Link access
 - Flow control
 - Error detection / correction
 - Half-duplex / full-duplex



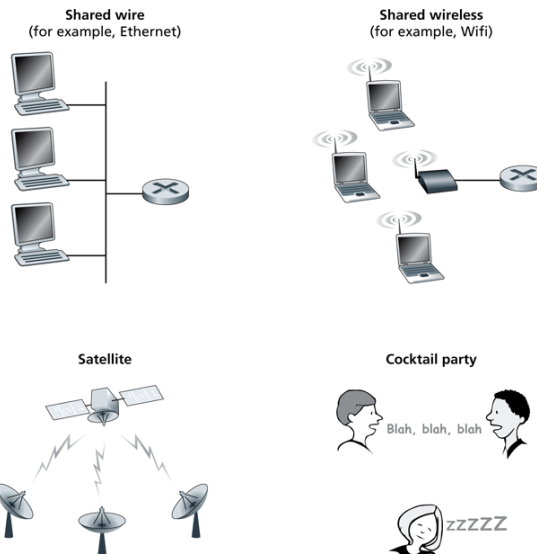
Network interface cards

- Implementation of NIC
 - Adapter connected to processor via bus
- NIC implements many functions in hardware
 - Error detection
 - Retransmission
 - Direct Memory Access (DMA)



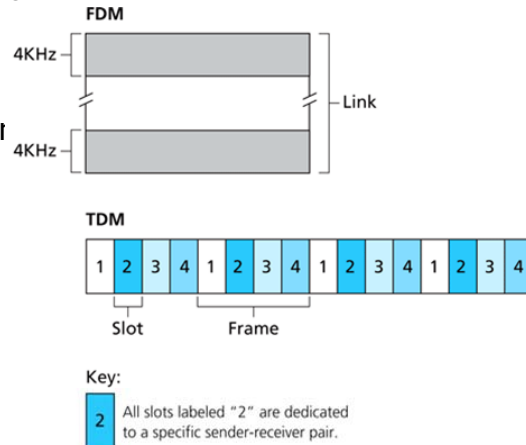
Multiple access

- Many data link layers use "shared medium"
 - Shared wire
 - Shared spectrum
- Medium Access Control (MAC) important
 - How can we coordinate sharing?



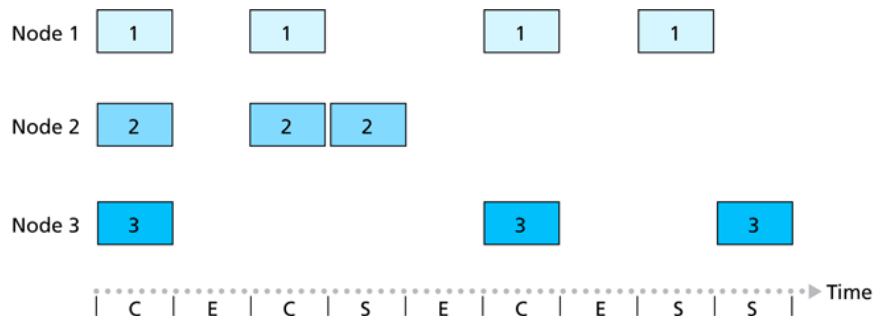
Multiple access

- Medium Access Control Approaches:
 - Channel partitioning
 - » TDM, FDM, CDMA
 - Random access protocols
 - » ALOHA, CSMA/CD
 - Taking turns protocol
 - » Polling, token-passing
- Random access used in Ethernet
 - How does it work?



Slotted ALOHA

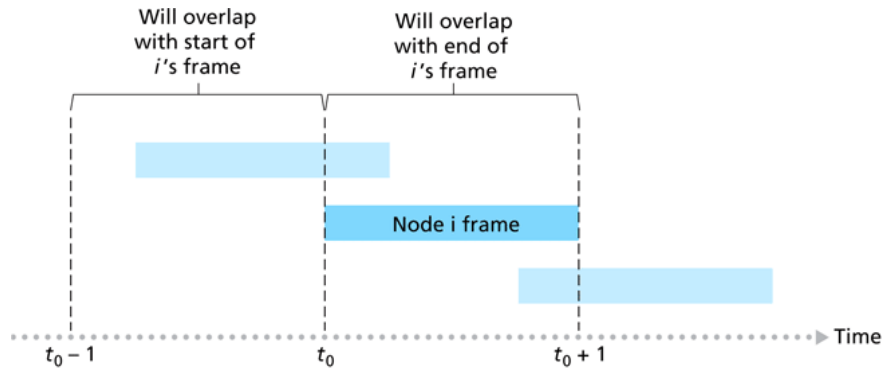
- Multiple stations may send at the same time?
 - What happens if collision occurs?



- Key:**
- C = Collision slot
 - E = Empty slot
 - S = Successful slot

Pure ALOHA

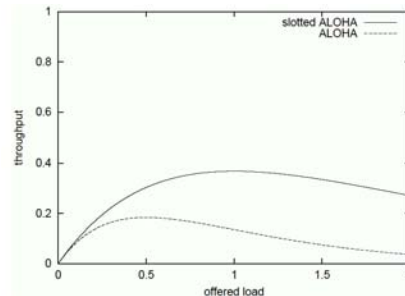
- Station can send whenever it is ready
 - Avoids synchronization challenge in Slotted ALOHA



- How does performance differ (pure vs. slotted)?

ALOHA throughput

- Slotted ALOHA analysis
 - Probability that k frames are generated during interval
 - Poisson distribution: $\Pr[k] = G^k e^{-G} / k!$
 - G is rate of frame generation (new and retransmission)
 - Probability of successful transmission
 - Success if no other frame generated (P_0): $S = GP_0 = Ge^{-G}$
- ALOHA analysis
 - Frame overlaps two slots
 - Probability that no other frame is generated during two slots:
 - $P_0 = e^{-2G}$
 - Success $S = GP_0 = Ge^{-2G}$
- Max throughput:
 - 37% @ $G = 1.0$ (slotted ALOHA)
 - 18% @ $G = 0.5$ (ALOHA)



ALOHA improvements

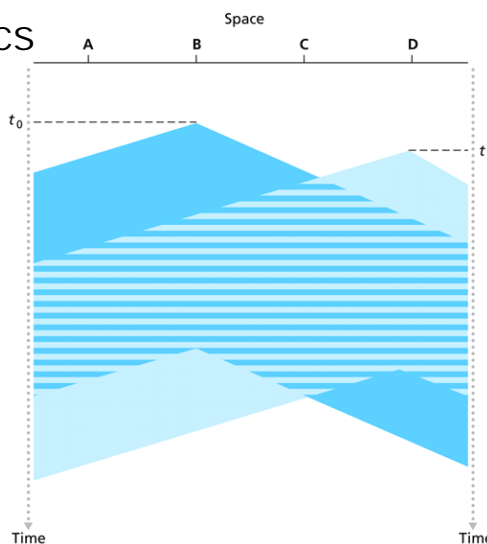
- How can we improve ALOHA?

ALOHA improvements

- How can we improve ALOHA?
 - Don't send when somebody else has already started
 - Stop when interference is already happening
 - » Why do we need to do that if we don't start sending when somebody else sends?
- Carrier Sensing (CS)
 - Listen on channel
 - Only send when nobody else is transmitting
- Collision Detection (CD)
 - Listen to own transmission on channel
 - If garbled then stop transmitting

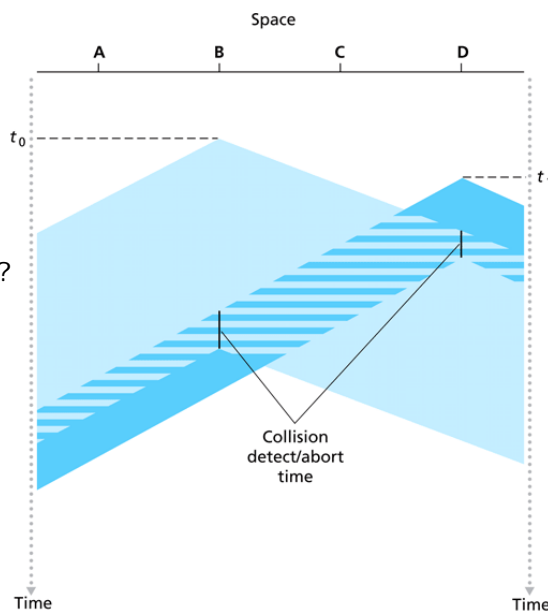
CSMA collision

- Collisions occur despite CS
 - Space-time diagram
- Propagation causes problems



CSMA/CD

- Collision detection reduces useless transmissions
- How can we guarantee that collision is detected?
 - Why is this necessary?



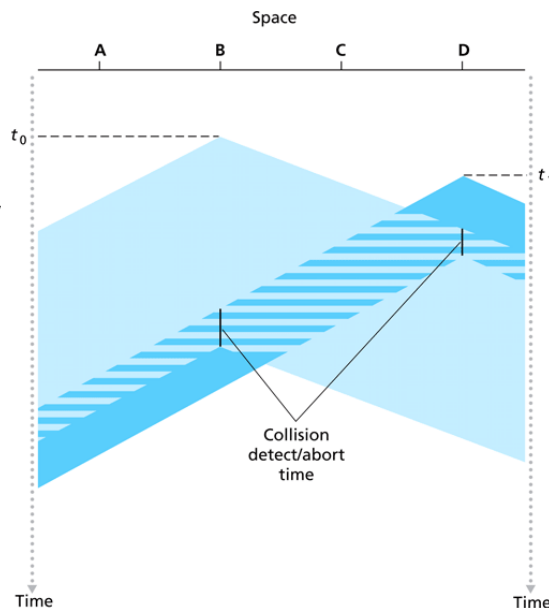
CSMA/CD

- How can we guarantee that collision is detected?

- Packet must be at least as long as twice the propagation delay between stations with maximum separation

- 10 Mbps Ethernet

- Packet length:
 - » 64 bytes
- Maximum distance:
 - » 2.5 km



Ethernet

- IEEE 802.3 protocol

- CSMA/CD

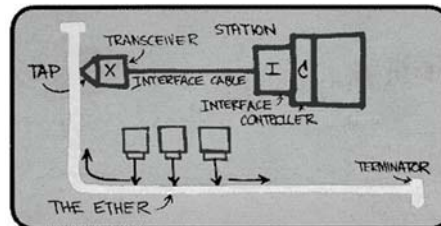
- Station listens and transmits when cable idle
- If collision occurs, then station backs off and retries

- Truncated exponential backoff

- Wait for random number of 512 bit-times
- After c collisions: uniform distribution over $[0 \dots 2^{\min\{c, 10\}} - 1]$
- After 16 collisions: transmission aborted

- Limitations

- Max 2.5km of cable, thus RTT limited to 51.2 μ s
- Corresponds to 64 bytes @ 10Mbps



Ethernet

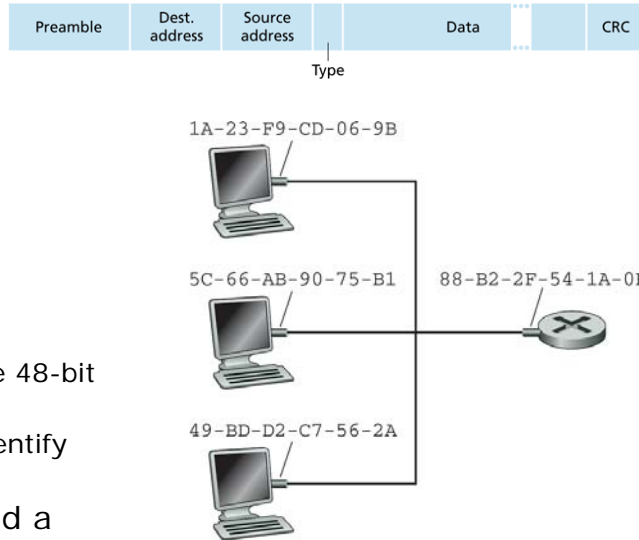
- Frame format:

- Preamble synchronized receiver
- Type identifies network layer protocol
- CRC for error detection

- Addressing

- Globally unique 48-bit address
- First 24 bits identify vendor

- How can we find a destination?



Address Resolution Protocol

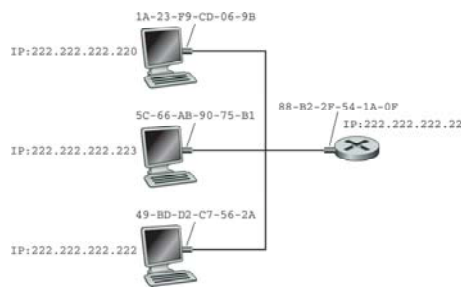
- Problem: what is Ethernet address of destination?

- ARP protocol

- Broadcast of request for Ethernet address of given IP address
- Response from anybody

- ARP table maintains entries

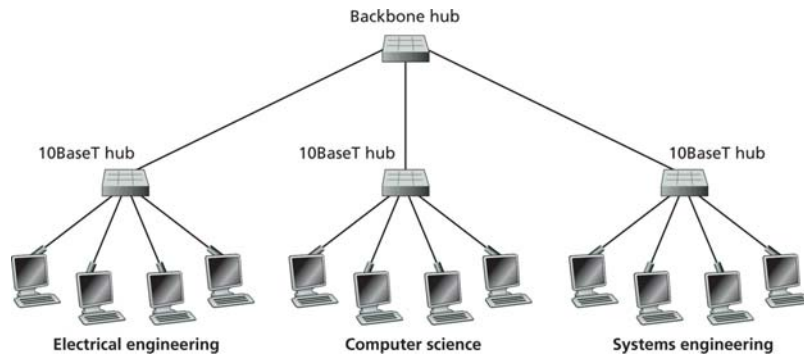
- Timeout ensures adaptation to reconfigurations



IP Address	MAC Address	TTL
222.222.222.221	88-B2-2F-54-1A-0F	13:45:00
222.222.222.223	5C-66-AB-90-75-B1	13:52:00

LAN Architecture

- LAN consists of several segments
- Hub connections maintain collision domain



ECE697AA – 09/16/08

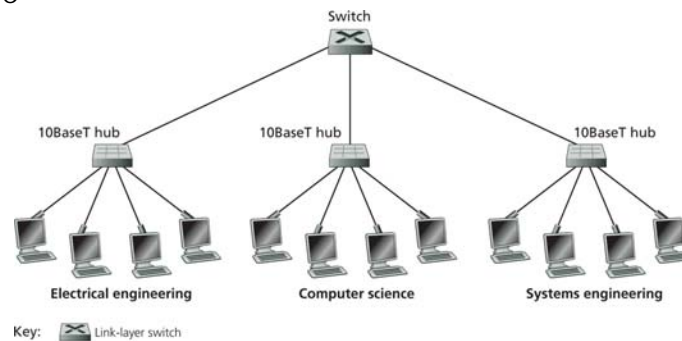
UMass Amherst – Tilman Wolf

17

LAN Architecture

- Switches separate collision domains
- Filtering and forwarding
 - Switch “knows” where frame should go
 - Switch table

Address	Interface	Time
01-12-23-34-45-56	2	9:39
62-FE-F7-11-89-A3	1	9:32
7C-BA-B2-B4-91-10	3	9:36
....



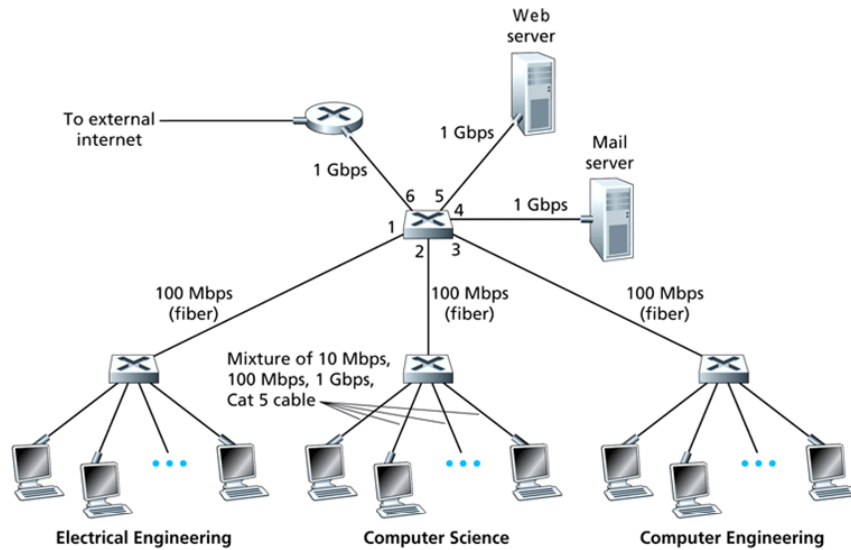
ECE697AA – 09/16/08

UMass Amherst – Tilman Wolf

18

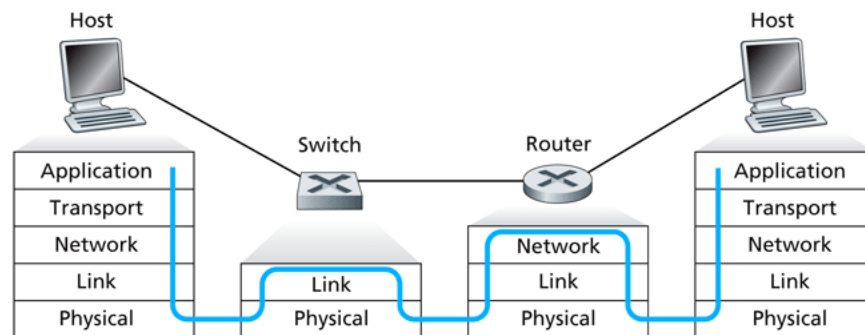
LAN Architecture

- Example with heterogeneous links



Network Architecture

- Comparison of LAN switches and routers



Assignments

- Read
 - kc claffy, Greg Miller, and Thompson Kevin, "The nature of the beast: Recent traffic measurements from an internet backbone," in Proc. of 1998 INET Conference, Geneva, Switzerland, June 1998.
- Have a look at
 - <http://www.caida.org>
- SPARK
 - Assessment quiz