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# **ECE 354 – Computer Systems Lab II**

Interrupts, Strings, and Busses

# Fun Fact

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- Press release from Microchip:  
*“Microchip Technology Inc. announced it provides PICmicro® field-programmable microcontrollers and system supervisors for the Segway Human Transporter (HT) [...] The **PIC16F87x Flash microcontrollers** process sensor data from the inertial monitoring unit and communicate information to the control module. Other **PIC16F87x** devices located in the battery packs provide monitoring functions. [...]”*



# Lab 1

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- All groups completed Lab 1 – good job!
  - Results posted on WebCT
- Understand how UART works, not just how to use it
- Questions?
- Additional lab hours – see web page
- Quiz: multiple answer possible

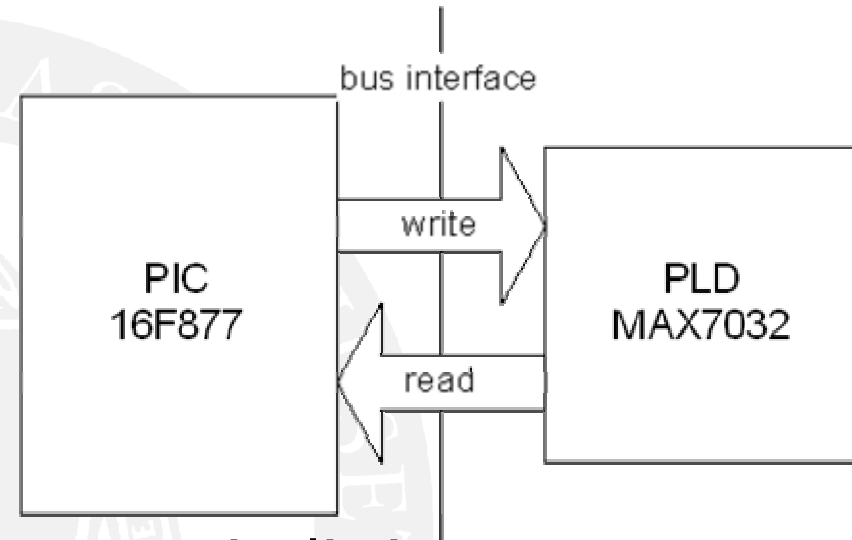
# Overview

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- Lab 2
- Interrupts
  - Why we need them
  - How to use them
- Timer
- Efficient printing of strings
- Interfacing PICs and PLDs
  - External bus design
  - READ and WRITE transactions

# Lab 2

- Interconnect PIC with PLD
- PLD acts as coprocessor
- PIC and PLD communicate via bus
- You have to design bus interface
- Timer on PIC is used to generate periodic interrupts
  - Will make sure your interface is robust
- PLD programmed in VHDL
  - Great example of hardware/software co-design!



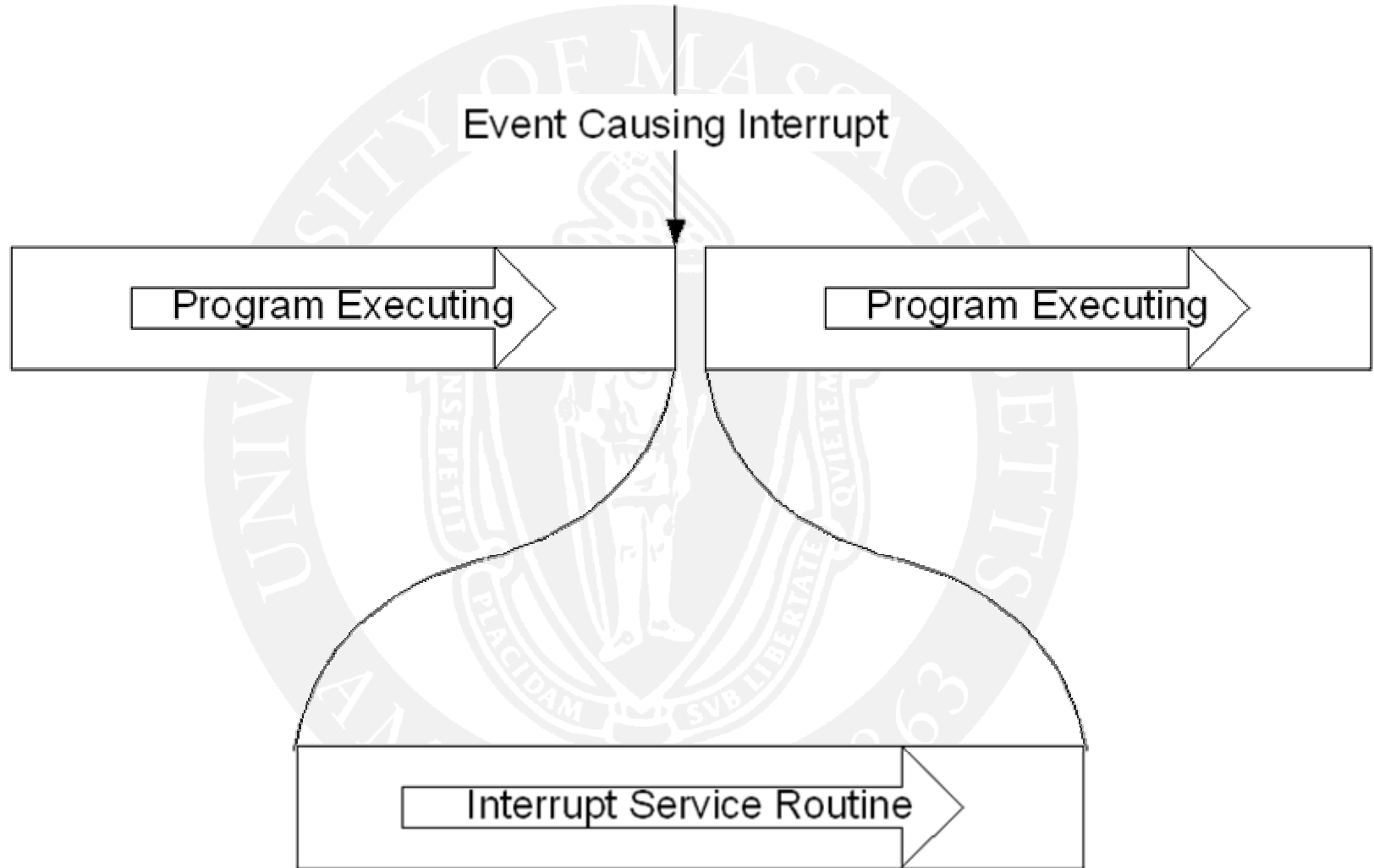
# Interrupts

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- Lab 1 used “polling”
  - What is bad about polling?
- Interrupts
  - Triggered by internal or external events
  - Cause program to “interrupt” and treat interrupt
  - After interrupt processing, processing returns to previous code
- What is better about interrupts?
- Example for interrupt triggers
  - UART transmission or reception
  - Change of input voltage on pin
  - Timer
  - A/D conversion completed

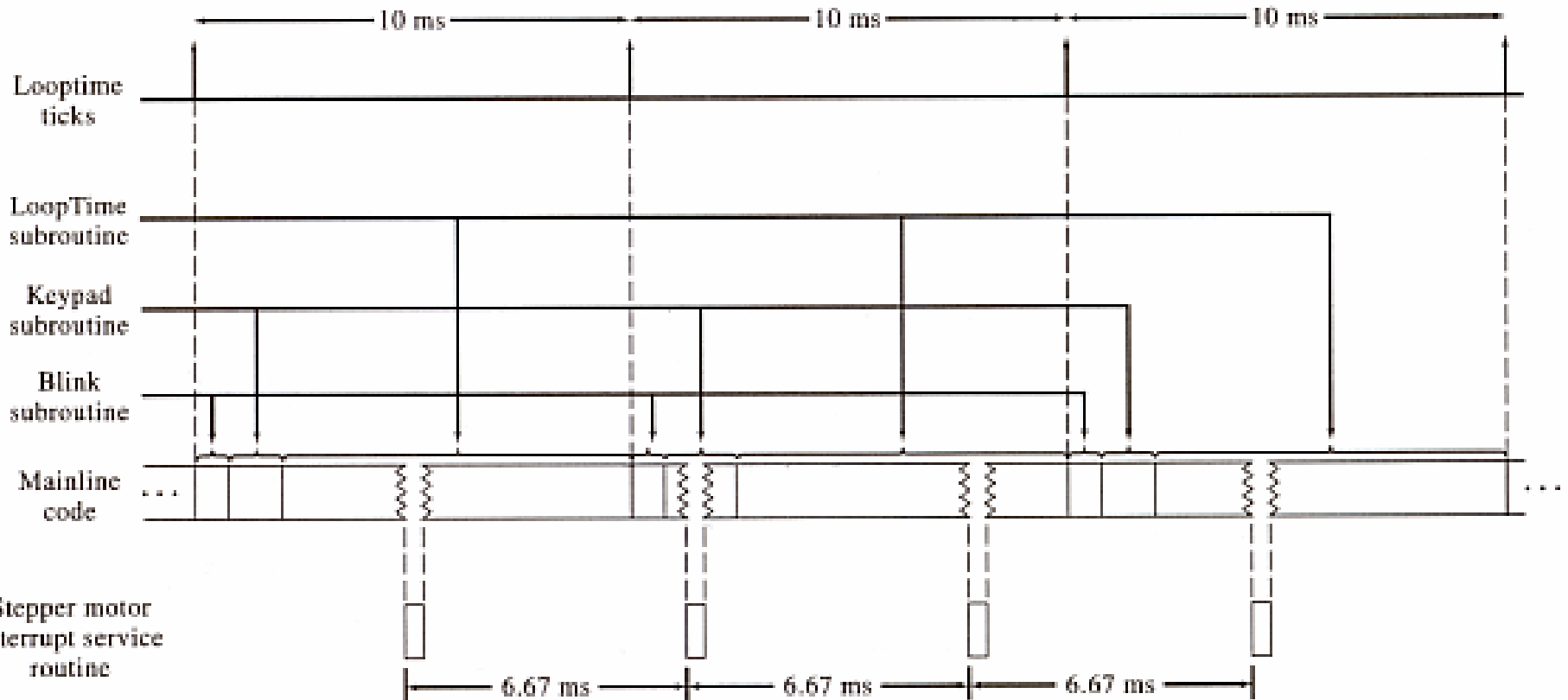
# Interrupt Concept

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# Interrupt Example

- More complex example:





# Steps During Interrupt

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1. Interrupts have to be enabled
  - Bits set in INTCON (internal interrupts) or PIE1 (peripheral interrupts) registers
2. Interrupt stimulus
  - Timer/counter overflows, change on external pin
3. Interrupts automatically disabled
  - Bit 7 of INTCON, why?
4. Jump to interrupt vector
  - Address 0x4, typically calls interrupt subroutine
5. Jump to and execute interrupt service routine
  - PIR1 register identifies which interrupt has triggered
6. Return to previous code
  - RETFIE instruction, also enables interrupts

# INTCON Register

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R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-x
GIE	PEIE	T0IE	INTE	RBIE	T0IF	INTF	RBIF
bit 7							bit 0

- bit 7      **GIE:** Global Interrupt Enable bit  
1 = Enables all unmasked interrupts  
0 = Disables all interrupts
- bit 6      **PEIE:** Peripheral Interrupt Enable bit  
1 = Enables all unmasked peripheral interrupts  
0 = Disables all peripheral interrupts
- bit 5      **T0IE:** TMR0 Overflow Interrupt Enable bit  
1 = Enables the TMR0 interrupt  
0 = Disables the TMR0 interrupt
- bit 4      **INTE:** RB0/INT External Interrupt Enable bit  
1 = Enables the RB0/INT external interrupt  
0 = Disables the RB0/INT external interrupt
- bit 3      **RBIE:** RB Port Change Interrupt Enable bit  
1 = Enables the RB port change interrupt  
0 = Disables the RB port change interrupt

# PIE1 Register

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
PSPIE <sup>(1)</sup>	ADIE	RCIE	TXIE	SSPIE	CCP1IE	TMR2IE	TMR1IE
bit 7							bit 0

bit 7 **PSPIE<sup>(1)</sup>**: Parallel Slave Port Read/Write Interrupt Enable bit

- 1 = Enables the PSP read/write interrupt
- 0 = Disables the PSP read/write interrupt

bit 6 **ADIE**: A/D Converter Interrupt Enable bit

- 1 = Enables the A/D converter interrupt
- 0 = Disables the A/D converter interrupt

bit 5 **RCIE**: USART Receive Interrupt Enable bit

- 1 = Enables the USART receive interrupt
- 0 = Disables the USART receive interrupt

bit 4 **TXIE**: USART Transmit Interrupt Enable bit

- 1 = Enables the USART transmit interrupt
- 0 = Disables the USART transmit interrupt

# PIR1 Register

R/W-0	R/W-0	R-0	R-0	R/W-0	R/W-0	R/W-0	R/W-0
PSPIF <sup>(1)</sup>	ADIF	RCIF	TXIF	SSPIF	CCP1IF	TMR2IF	TMR1IF
bit 7							bit 0

- bit 7 **PSPIF<sup>(1)</sup>**: Parallel Slave Port Read/Write Interrupt Flag bit  
1 = A read or a write operation has taken place (must be cleared in software)  
0 = No read or write has occurred
- bit 6 **ADIF**: A/D Converter Interrupt Flag bit  
1 = An A/D conversion completed  
0 = The A/D conversion is not complete
- bit 5 **RCIF**: USART Receive Interrupt Flag bit  
1 = The USART receive buffer is full  
0 = The USART receive buffer is empty
- bit 4 **TXIF**: USART Transmit Interrupt Flag bit  
1 = The USART transmit buffer is empty  
0 = The USART transmit buffer is full

# RETFIE Instruction

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RETFIE                      Return from Interrupt

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Syntax:                      [*label*] RETFIE

Operands:                    None

Operation:                   TOS → PC,  
                                  1 → GIE

Status Affected:          None



# Interrupt Code Sample

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```
org    H'000'    ; Reset vector
goto   Mainline; Location of start of program

org    H'004'    ; Interrupt vector
goto   IntServ  ; Start of int service routine

Mainline
....
....

org    H'100'    ; put service routing at 0x100
IntServ
....           ; first inst. of service routine
....
retfie        ; return from interrupt instr.
```

# Saving State

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- Some “state” of PIC is not preserved during interrupt
  - What is “state”?
  - What is preserved?
  - What can get lost?
- How to avoid problems:
  - Preserve state (w, STATUS) before interrupt processing or
  - Do not change state during interrupt processing (difficult)
- Moving w and STATUS to temporary variables does not help! Why?
  - movf causes “evaluation” of value and impacts STATUS
  - Note: use of swapf instruction instead of movf
- See section 12.11 of data sheet and Peatman section 4.5

# Interrupt Limitations

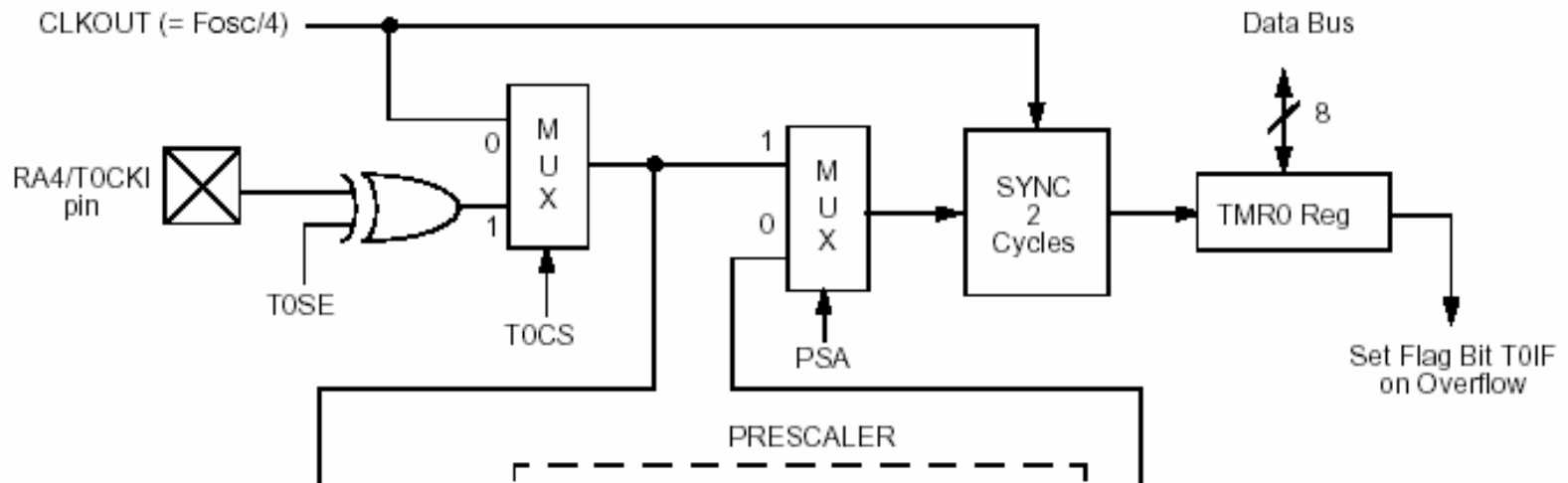
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- What happens if interrupt is too long?
  - Other critical interrupts cannot be handled or
  - Livelock (not on PIC due to lack of recursion)
- Beware of function calls in interrupt service routine
  - Stack overflow could happen
  - $\text{max nesting of program} + \text{max nesting of ISR} + 1 \leq 8$



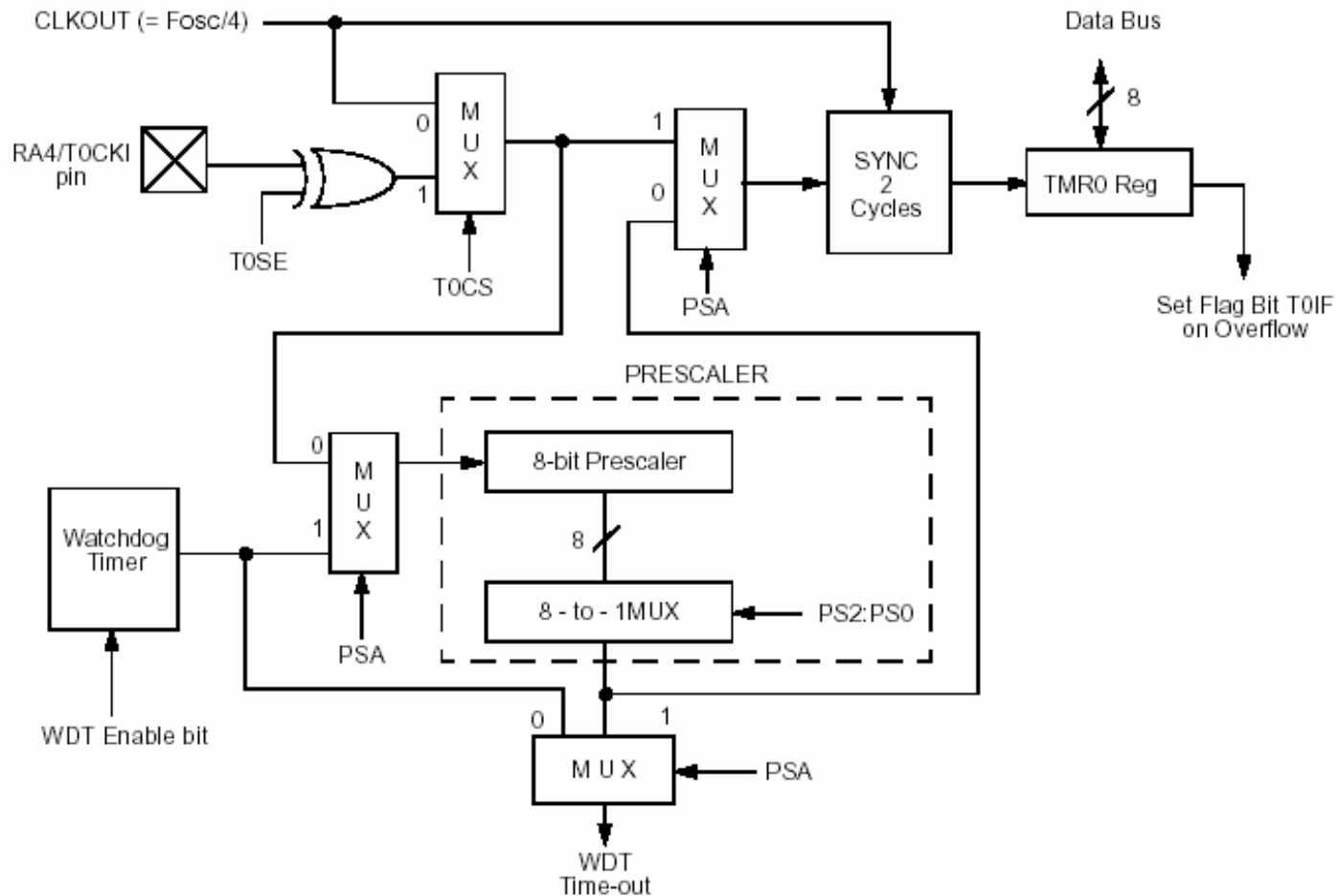
# Timer Interrupts

- The PIC has several built-in timers
- timer0 is a simple 8-bit counter
  - External or internal clock
  - Prescaler possible
- See Peatman pp.100–103, data sheet pp.47–49



# Prescaler

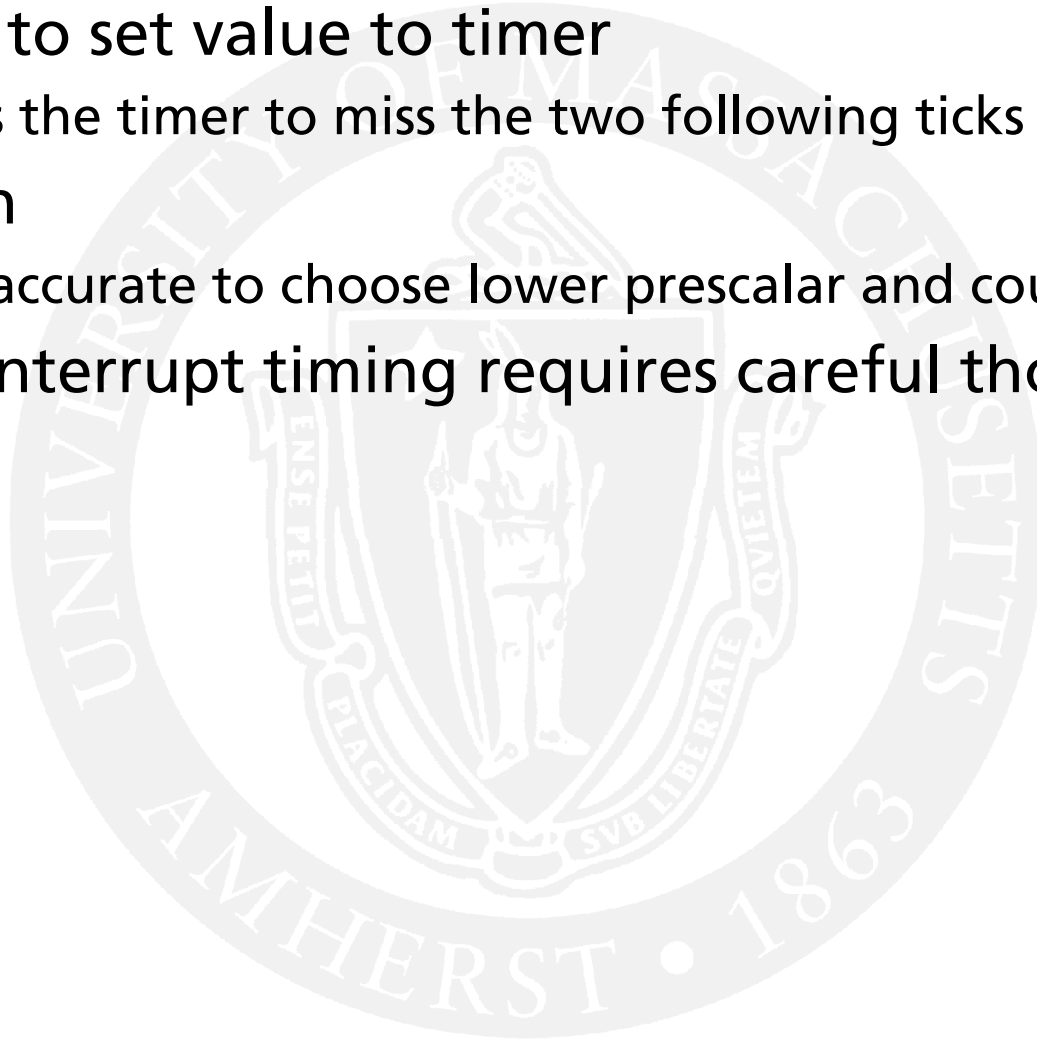
- Prescaler performs additional counting. Why bother?



# Timer Details

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- Possible to set value to timer
  - Causes the timer to miss the two following ticks
- Precision
  - More accurate to choose lower prescaler and count more
- Precise interrupt timing requires careful thought



# String Printing

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- Printing strings on terminal can be awkward
  - Example: print "Hello"
    - call wait ; subroutine which checks PIR bit 4.
    - movlw 'H' ; send ASCII w char. to W
    - movwf TXREG ; send w char. to UART trans. buffer
    - call wait ; subroutine which checks PIR bit 4.
    - movlw 'e' ; send ASCII o char. to W
    - movwf TXREG ; send o char. to UART trans. buffer
    - call wait ; subroutine which checks PIR bit 4.
    - movlw 'l' ; send ASCII w char. to W
    - movwf TXREG ; send w char. to UART trans. Buffer
    - ...
- Any ideas?

# Advanced String Handling (1)

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- Store text into (program!) memory
- Process one character a time
- Consider:

```
    ORG    0x0
    goto   Start ; jump to the start of the program
    ORG    0x5
sub1  nop      ; start of a subroutine
    ...
    return   ; return from subroutine
    ORG    0x30
Start
    ...
    call   sub1 ; call the routine
    nop    ; first instr. after return
```

# Advanced String Handling (2)

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- call instruction:

CALL	Call Subroutine
Syntax:	[ <i>label</i> ] CALL k
Operands:	$0 \leq k \leq 2047$
Operation:	(PC)+ 1 → TOS, k → PC<10:0>, (PCLATH<4:3>) → PC<12:11>

- return instruction:

RETURN	Return from Subroutine
Syntax:	[ <i>label</i> ] RETURN
Operands:	None
Operation:	TOS → PC
Status Affected:	None

# Advanced String Handling (3)

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- retlw instruction returns and puts value into w

RETLW                      Return with Literal in W

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Syntax:                    [*label*] RETLW *k*

Operands:                 $0 \leq k \leq 255$

Operation:                 $k \rightarrow (W)$ ;  
                               $TOS \rightarrow PC$

- Use retlw for string printing
  - Call subroutine for each character
  - Use retlw to return each character and place into w
  - Send w to UART

# String Handling Code (1)

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```
BANKORAM    EQU      H'20'      ; equate a constant to hex 20.
            ORG      BANKORAM   ; reserve space in DATA MEMORY
            cblock    ; create a pointer in bank 0 at 0x20
            POINTER   ; name of value
            endc

            ORG      0x0
            goto     Start      ; jump to the start of the program

            ORG      0x5
sub1        movf     POINTER, W  ; move value in POINTER to W
            addwf    PCL, F      ; add value to PC
            retlw    A'H'
            retlw    A'e'
            retlw    A'l'
            retlw    A'l'
            retlw    A'o'
            retlw    0
            RETURN              ; shouldn't get here
```



# String Handling Code (2)

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```

Start
Loop
    ORG      0x30
    clr     POINTER
    call    sub1
    ...
    ; check if return value is 0
    btfsc  status, z      ; branch if not 0
    goto   Done          ; else done
    ...
    ; check bit 4 in PIR
    movwf  TXREG
    ...
    ; increment POINTER
    goto   Loop          ; print another character
    ORG    0x60

Done
    nop
    goto  Done
```

# String Handling

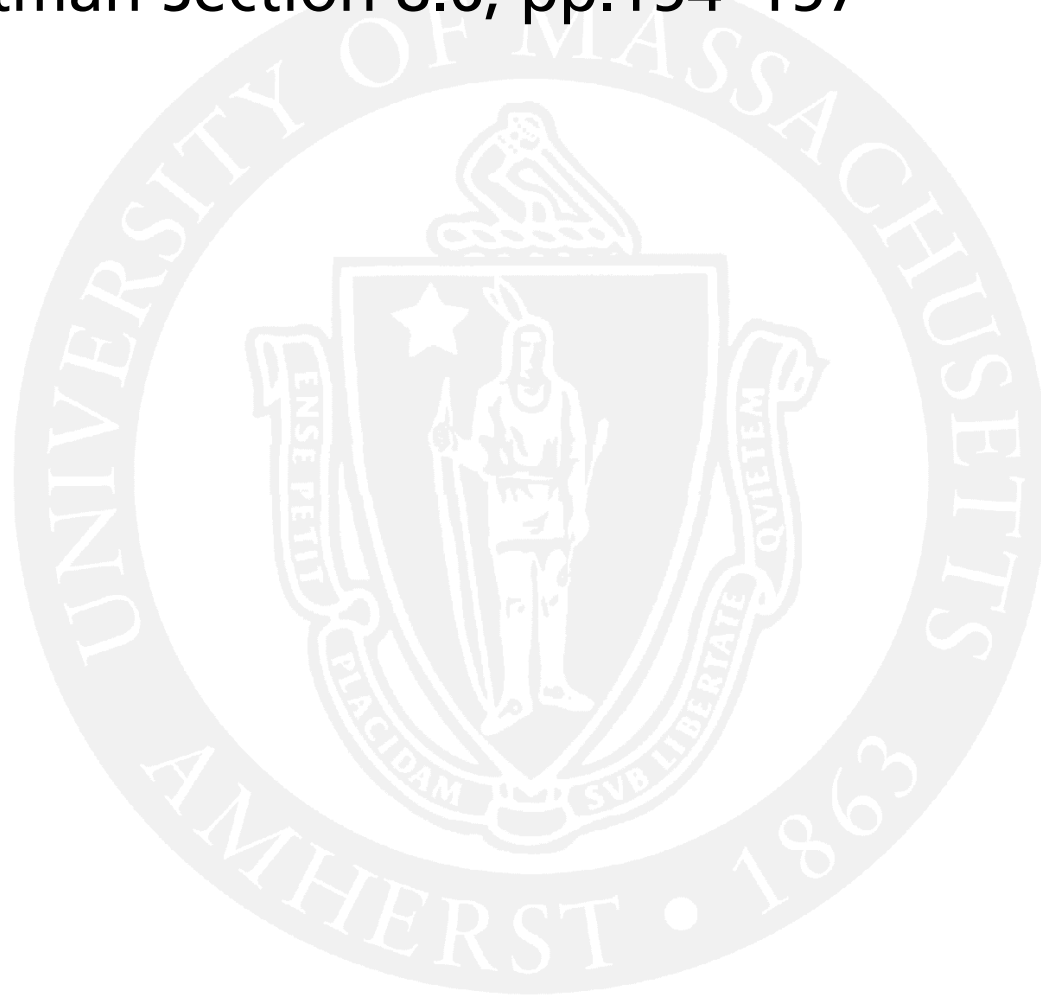
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- A bit complex
  - Simplified by 'dt' assembler directive  
dt "Hello" translates into  
retlw A'H'  
retlw A'e'  
retlw A'l'  
retlw A'l'  
retlw A'o'
- Note: does not terminate string (with A'0')!
- Saves memory space
- More easily modifiable
- You need to understand how it works, but you don't need to use it.

# OK, Who's Confused?

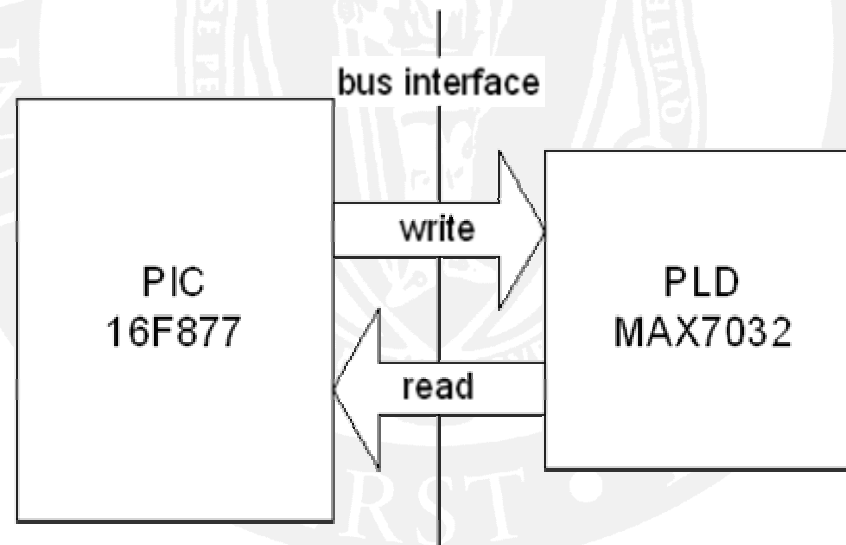
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- See Peatman Section 8.6, pp.154–157



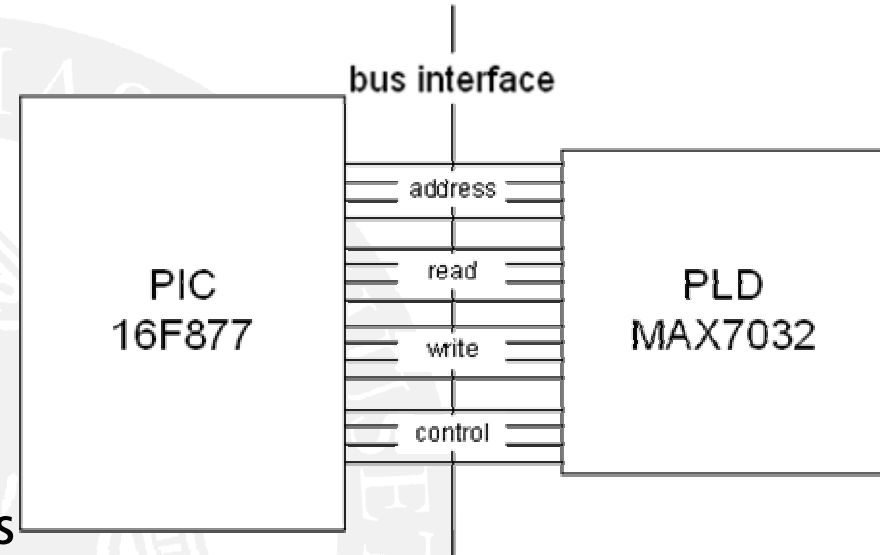
# Data Exchange Via Bus

- Multiple devices can be connected through a bus
  - We connect PIC to PLD
- PIC should be able to read from and write to PLD
  - PLD acts as coprocessor
  - For example: write value to 0x1 and read value+1 from 0x2
  - Functions: increment, bit count, maximum (since startup)



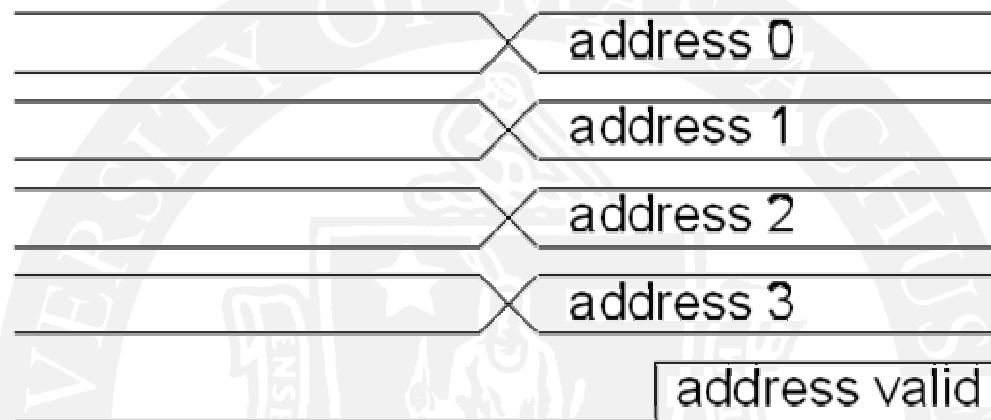
# Bus Interface

- A bus needs:
  - Address, data, control signals
- For Lab 2:
  - Port A: four bit address value
  - Port B: four bit data input
  - Port D: four bit data output
  - Port C: up to six control signals
- Build “read” and “write” transactions
  - Need to be robust and reliable (interrupts!)



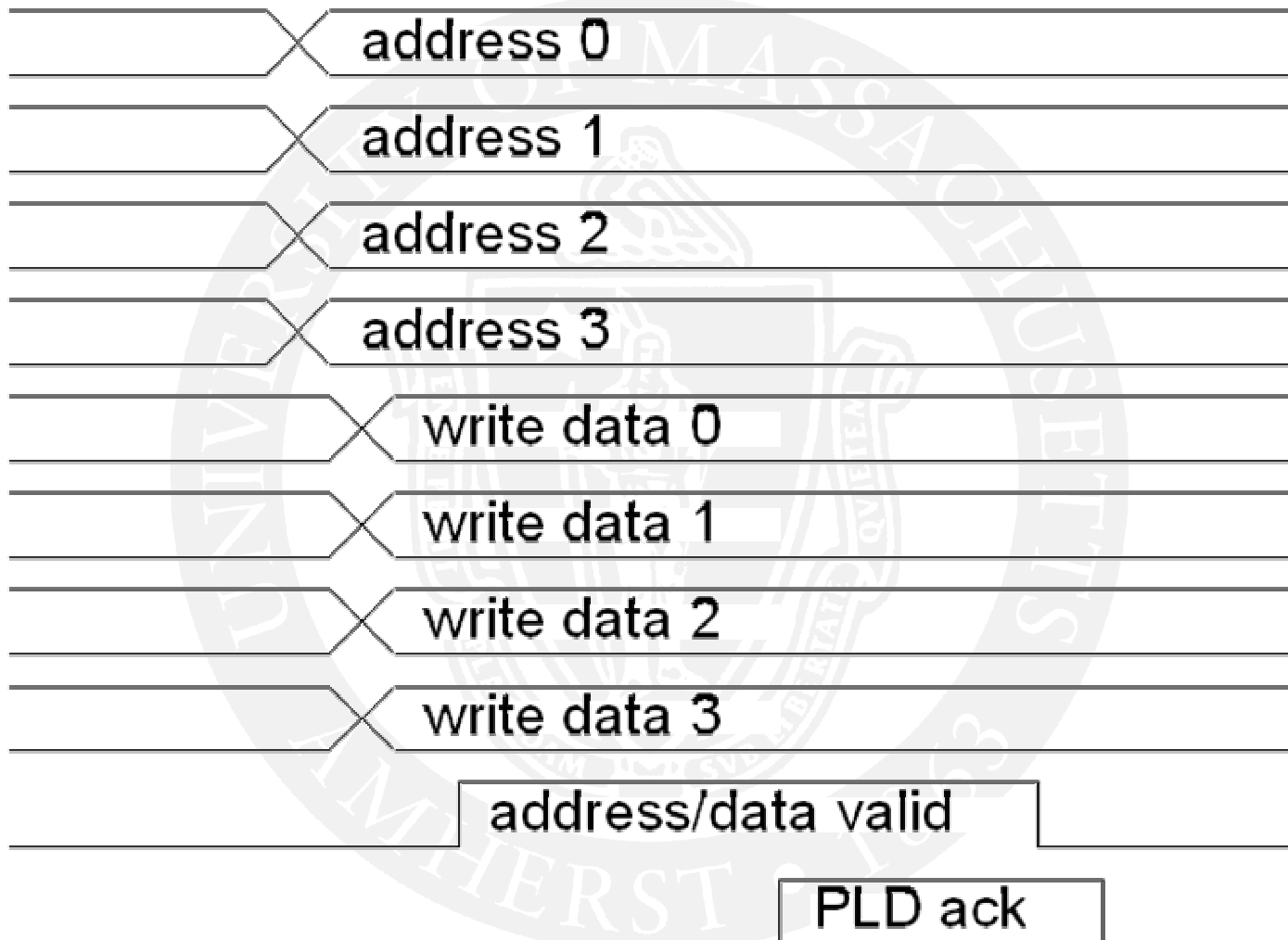
# Bus Issues

- Control signals indicate valid address/data

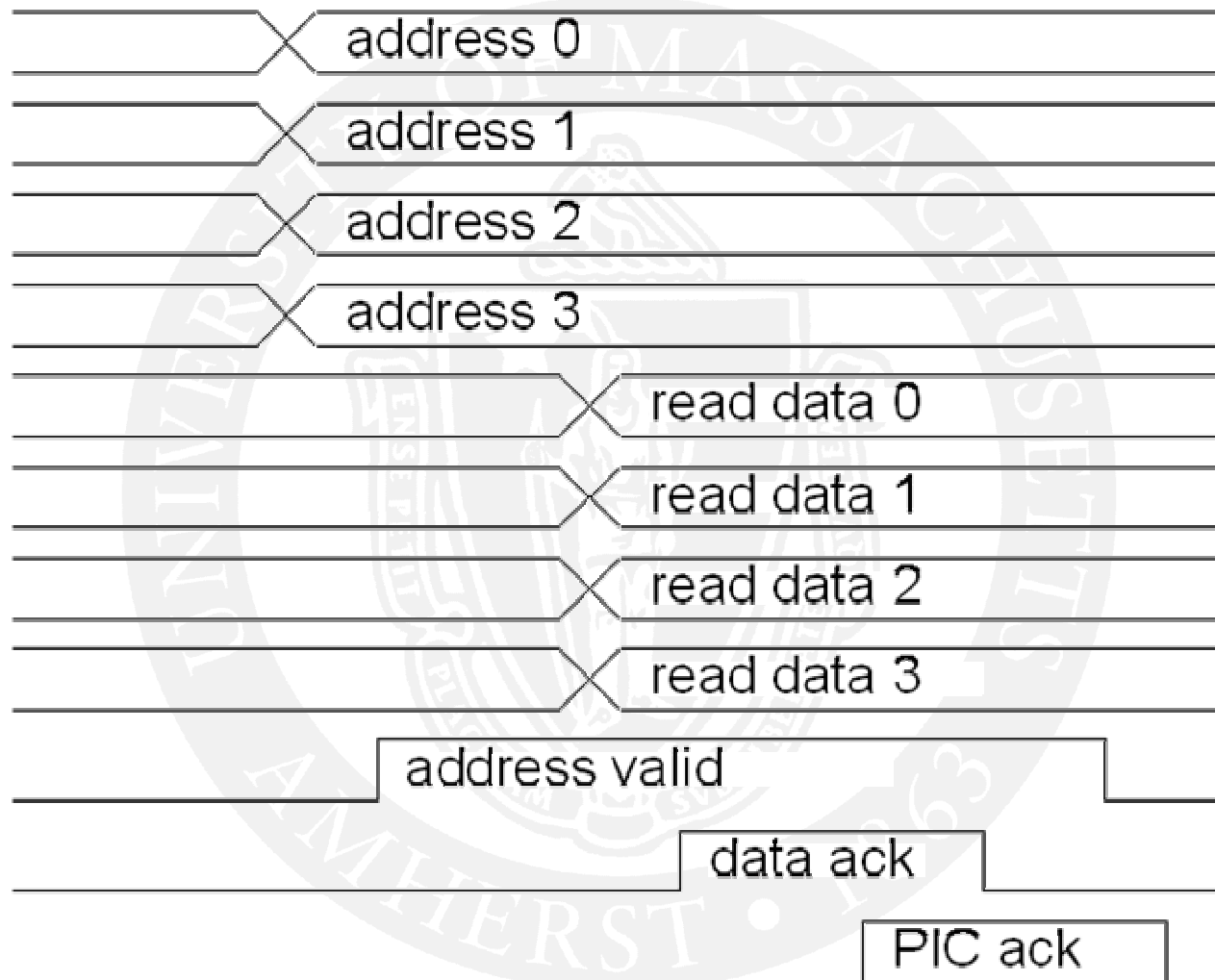


- PIC and PLD use same clock (synchronous)
  - Interrupts on PIC can cause delays!
- Transactions needs to be acknowledged (why?)
  - PLD acks when WRITE result was received
  - PIC acks when READ result was received
- The logic analyzer is your best friend 😊

# Example WRITE



# Example READ





# Bus Implementation

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- Need signal to distinguish READ and WRITE
- There are better, simpler ways
  - Consider using control bits for multiple purposes
- Use PLD state machines
  - One for address
  - One for read
  - One for write
- Implement on PLD using VHDL
  - Basically what you have done in ECE 353

# Lab 2

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- **BEFORE YOU START: READ!**
  - Lab Assignment
  - PIC data sheet section 12.10 – 12.11 (interrupts)
  - Peatman, chapter 4 (timer and interrupts)
- **Think about how you want to split work**
  - If you separate PIC and PLD design, make sure you have a good bus interface!
- **Think about steps to take to get it working**
- **Start working early!**
  - You will need more time than for Lab 1
  - Quiz March 1<sup>st</sup> – 3<sup>rd</sup>
  - Lab demos March 25<sup>th</sup> and 26<sup>th</sup>