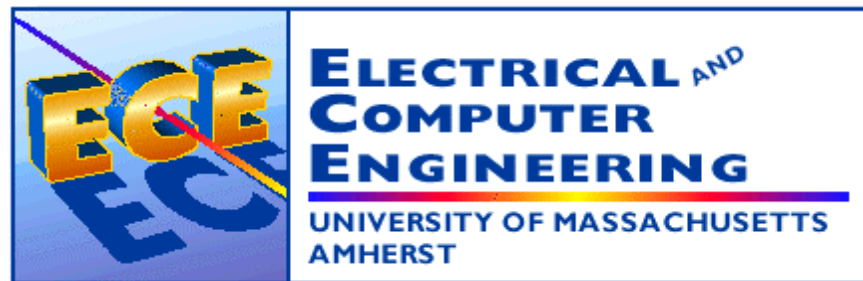

ENGIN 112

Intro to Electrical and Computer Engineering

Lecture 33

Arithmetic Logic Unit (ALU)

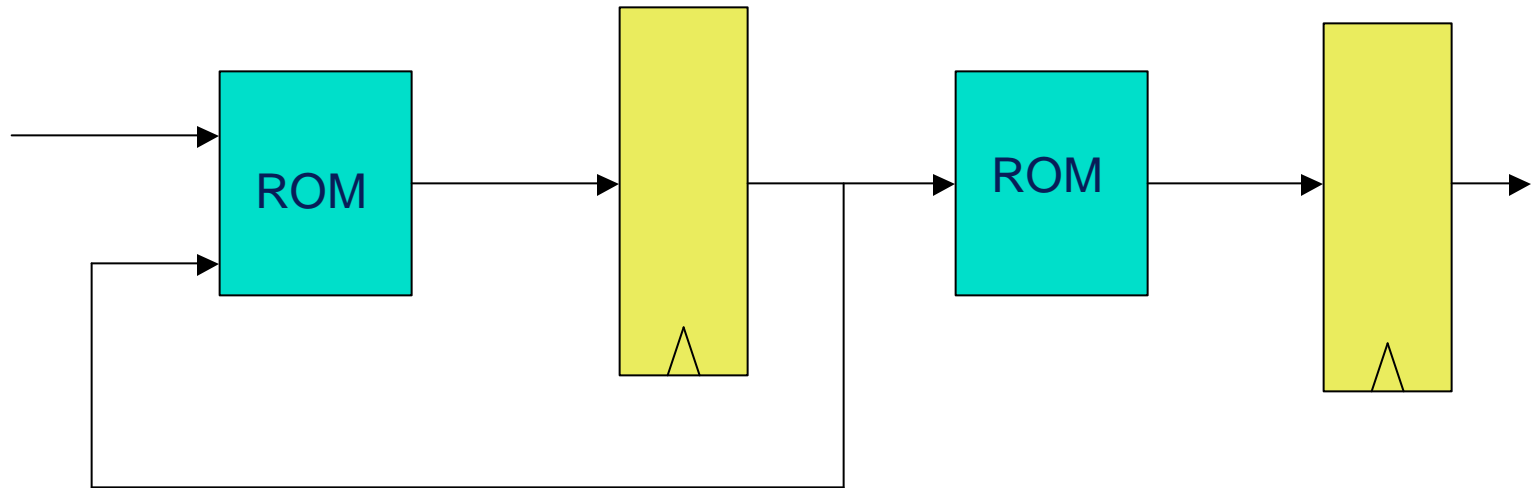


Overview

- **Main computation unit in most computer systems**
- **ALUs perform a variety of different functions**
 - Add, subtract, OR, AND...
- **Example: ALU chip (74LS382)**
 - Has data and control inputs
- **Individual chips can be chained together to make larger ALUs**
- **ALUs are important parts of datapaths**
 - ROMs often are used in the control path
- **Build a data and control path**

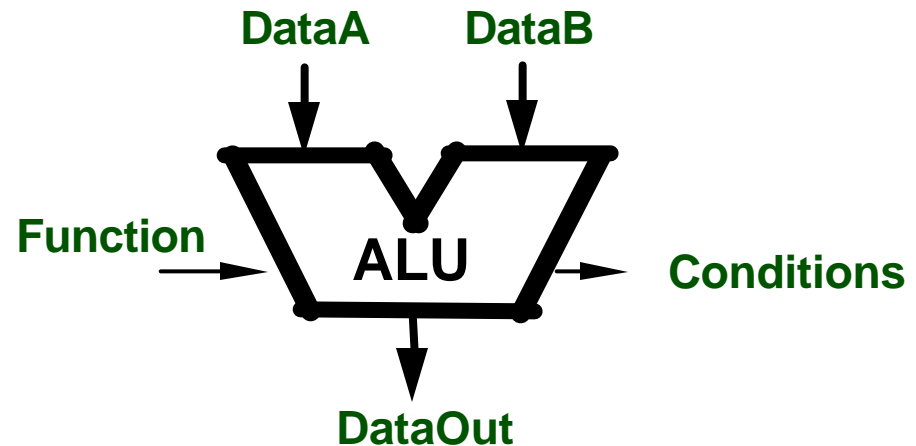
ROM-based Moore Machine Timing

- What is the maximum clock frequency of this circuit?
- Does this circuit satisfy **hold time** constraints?

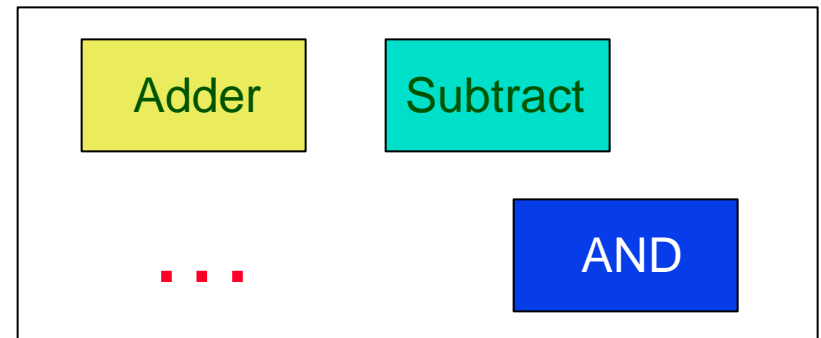


Arithmetic Logic Unit

- Arithmetic logic unit functions
 - Two multi-bit data inputs
 - **Function** indicates action (e.g. add, subtract, OR...)
- **DataOut** is same bit width as multi-bit inputs (**DataA** and **DataB**)
- ALU is combinational
- **Conditions** indicate special conditions of arithmetic activity (e.g. overflow).

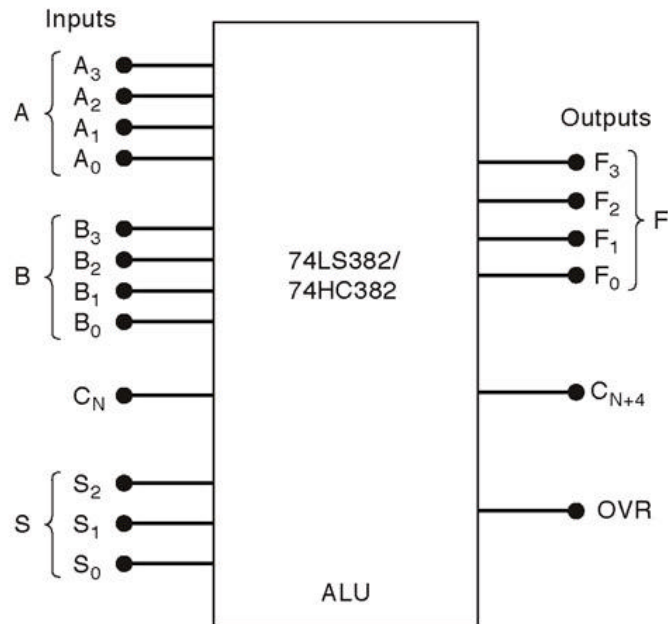


Think of ALU as a number of other arithmetic and logic blocks in a single box! **Function** selects the block



ALU Integrated Circuit

- Integrated circuit – off-the-shelf components
- Examine the functionality of this ALU chip



A = 4-bit input number F = 4-bit output number
 B = 4-bit input number C_{N+4} = carry out of MSB position
 C_N = carry into LSB position OVR = overflow indicator
 S = 3-bit operation select inputs

(a)

Function Table

S_2	S_1	S_0	Operation	Comments
0	0	0	CLEAR	$F_3F_2F_1F_0 = 0000$
0	0	1	B minus A	} Needs $C_N = 1$
0	1	0	A minus B	
0	1	1	A plus B	Needs $C_N = 0$
1	0	0	$A \oplus B$	Exclusive-OR
1	0	1	$A + B$	OR
1	1	0	AB	AND
1	1	1	PRESET	$F_3F_2F_1F_0 = 1111$

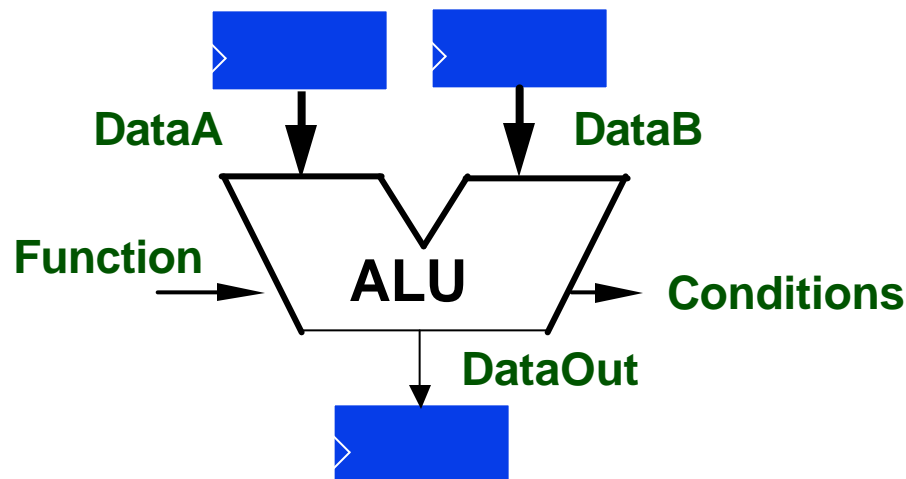
Notes: S inputs select operation.
 OVR = 1 for signed-number overflow.

(b)

Performs 8 functions

Example

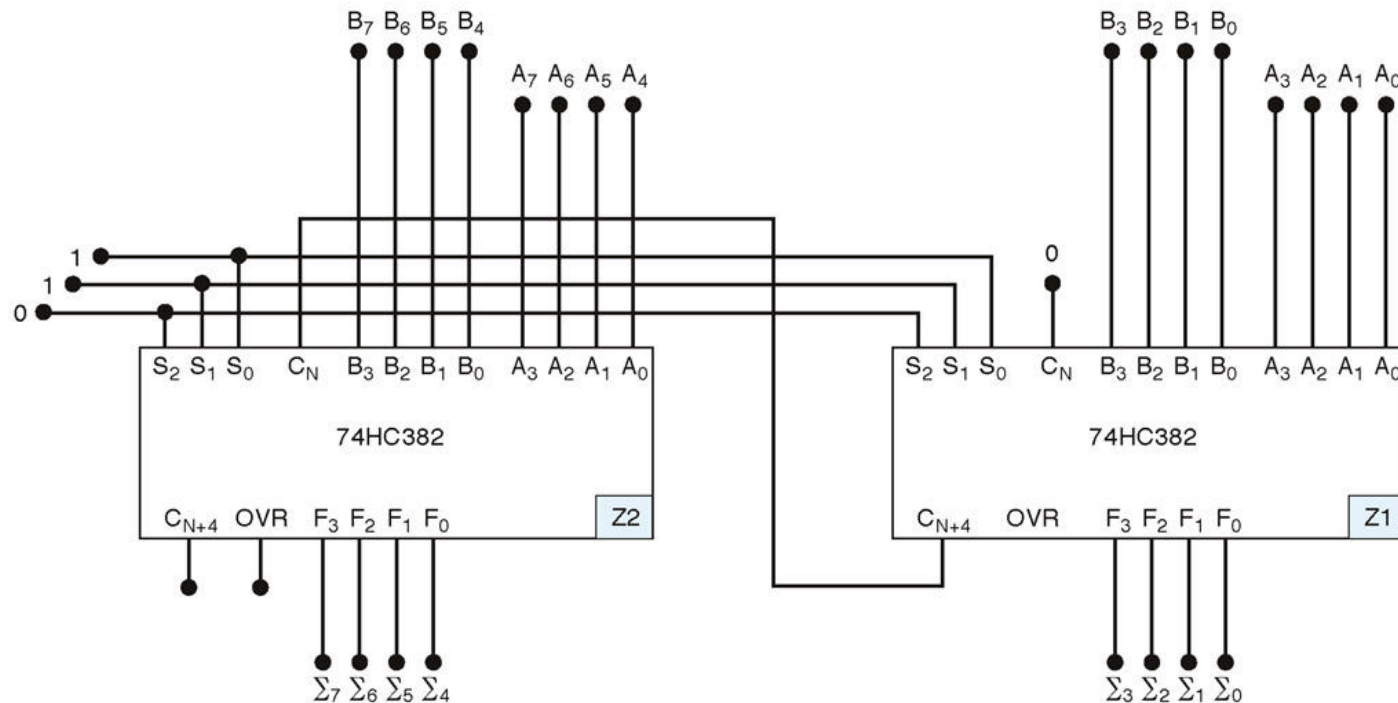
- Determine the 74HC382 ALU outputs for the following inputs: $S_2S_1S_0=010$, $A_3A_2A_1A_0=0100$, $B_3B_2B_1B_0=0001$, and $C_N=1$.
 - Function code indicates **subtract**
 - **$0100 - 0001 = 0011$**
- Change the select code to 101 and repeat.
 - Function code indicates **OR**
 - **$0100 \text{ OR } 0001 = 0101$**



Synchronize ALU
with a clock

Expanding the ALU

- Multi-bit ALU created by connecting carry output of low-order chip to carry in of high order

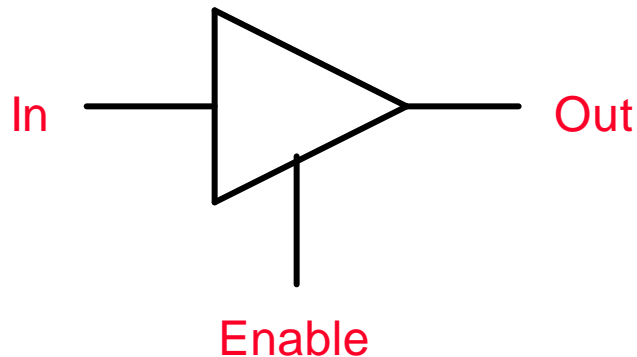


Notes: Z1 adds lower-order bits.
Z2 adds higher-order bits.
Σ₇–Σ₀ = 8-bit sum.
OVR of Z2 is 8-bit overflow indicator.

Eight-bit ALU formed from 2 four-bit ALUs

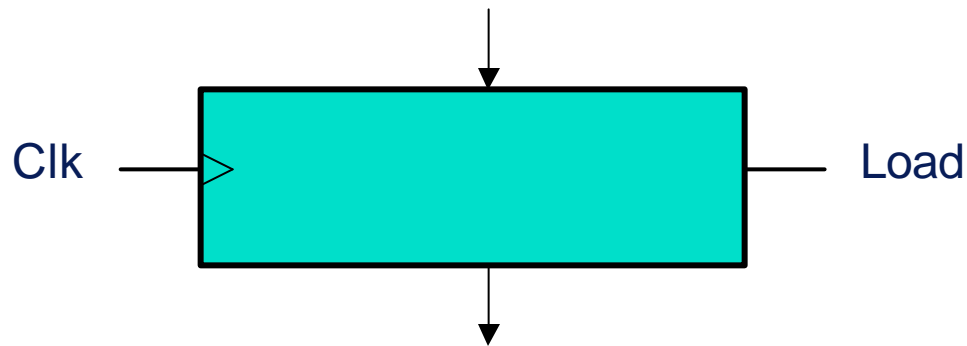
Datapath components

◦ Tri-state buffer



If Enable asserted,
Out = In
Otherwise
Out open-circuit

◦ Loadable register



Data stored on rising edge if Load is asserted (e.g. Load = 1)

Computation in a Typical Computer

- Control logic often implemented as a finite state machine (including ROMs)
- Datapath contains blocks such as ALUs, registers, tri-state buffers, and RAMs
- In a processor chip often a 5 to 1 ratio of datapath to control logic

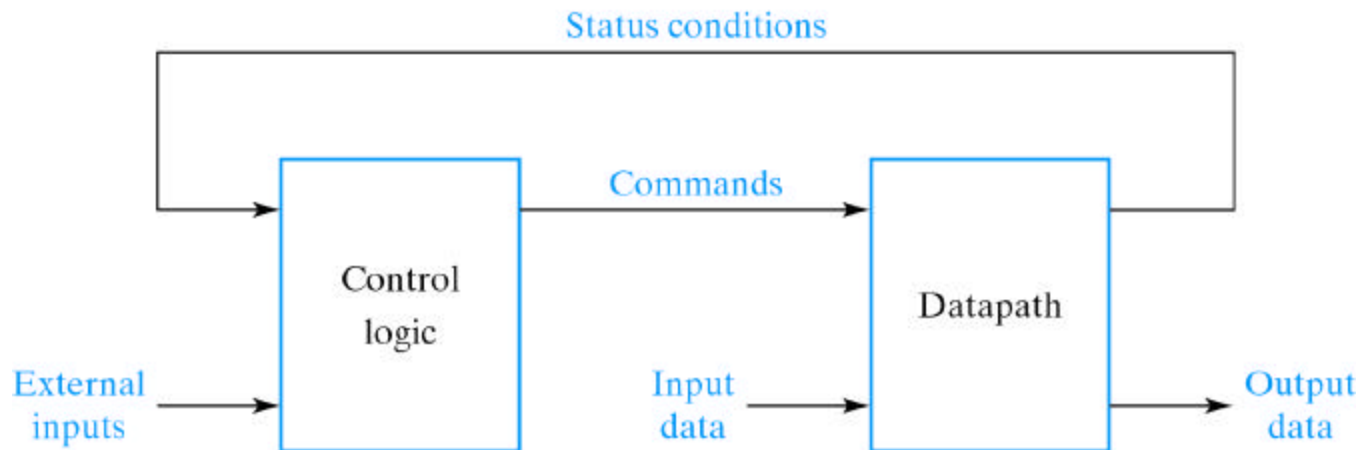


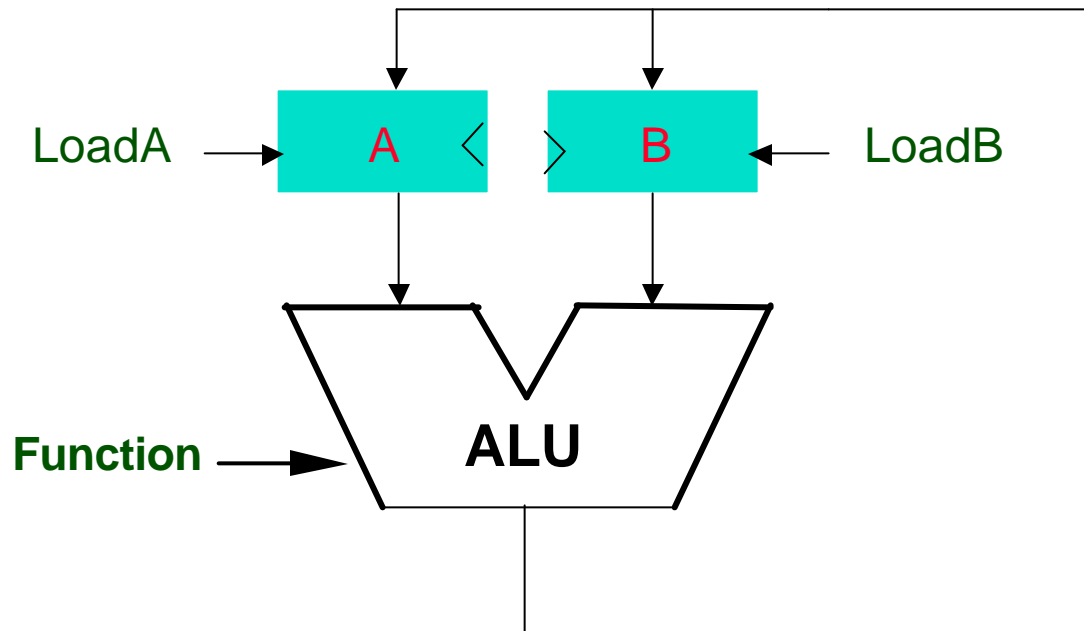
Fig. 8-2 Control and Datapath Interaction

Using a Datapath

Consider the following computation steps

1. ADD A, B and put result in A
2. Subtract A, B and put result in B
3. OR A, B put result in A
- Repeat starting from step 1

Determine values
for Function, LoadA, LoadB



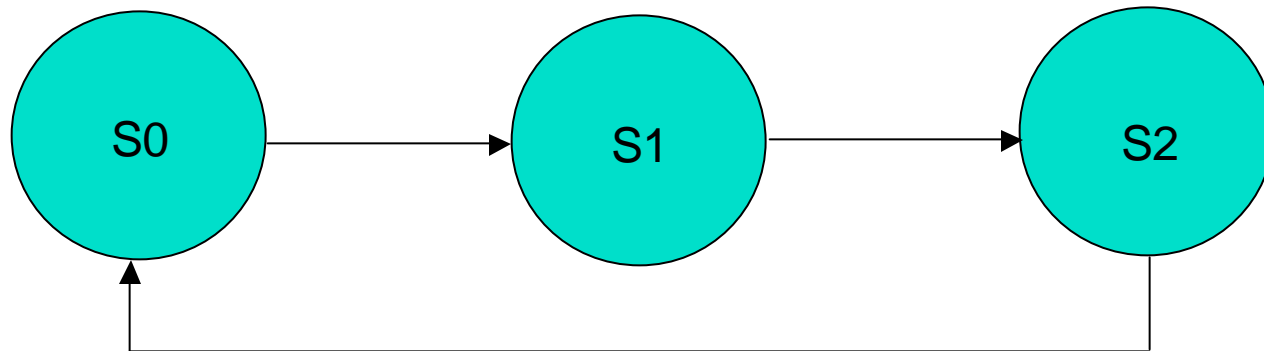
Modeling Control as a State Machine

◦ **Consider the following computation steps**

1. **ADD A, B and put result in A**
2. **Subtract A, B and put result in B**
3. **OR A, B put result in A**
- **Repeat starting from step 1**

Determine values
for Function, LoadA, LoadB

Model control as a state machine.
Determine control outputs for each state



Modeling Control as a State Machine

◦ **Consider the following computation steps**

1. **ADD A, B and put result in A**
2. **Subtract A, B and put result in B**
3. **OR A, B put result in A**
 - **Repeat starting from step 1**

States

S0 = 00

S1 = 01

S2 = 10

Present State	Next State	Function	LoadA	LoadB
00	01	011	1	0
01	10	010	0	1
10	00	101	1	0

We know how to implement this using an SOP.
Can we use a ROM?

ROM Implementation of State Machine

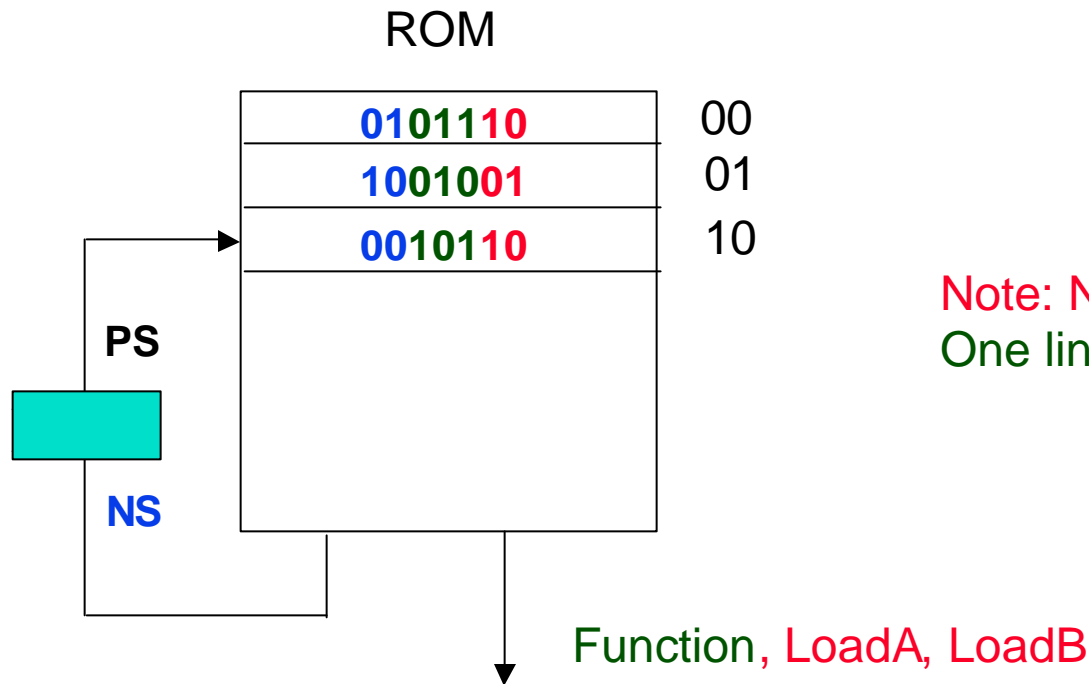
Present State	Next State	Function	LoadA	LoadB
00	01	011	1	0
01	10	010	0	1
10	00	101	1	0

States

S0 = 00

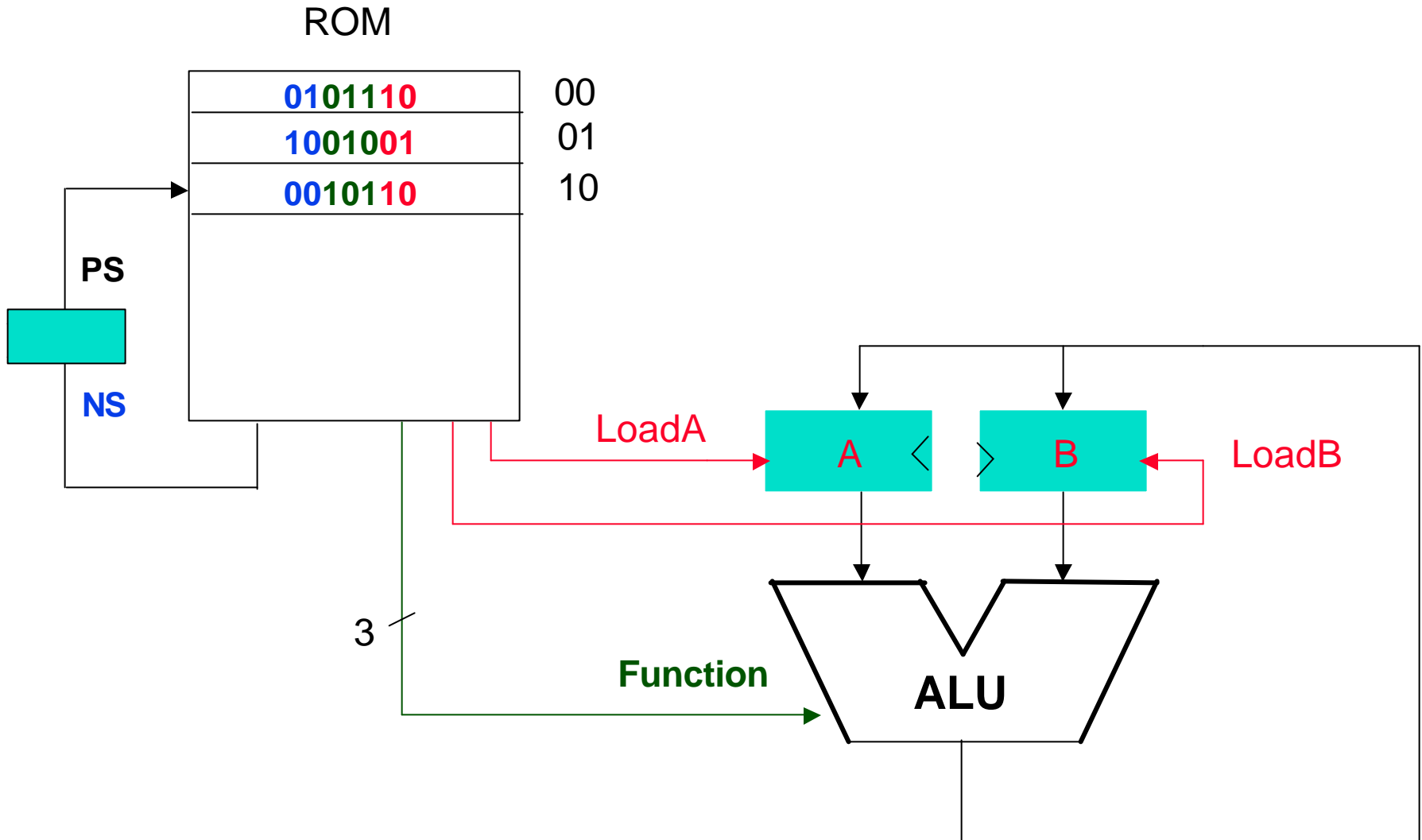
S1 = 01

S2 = 10

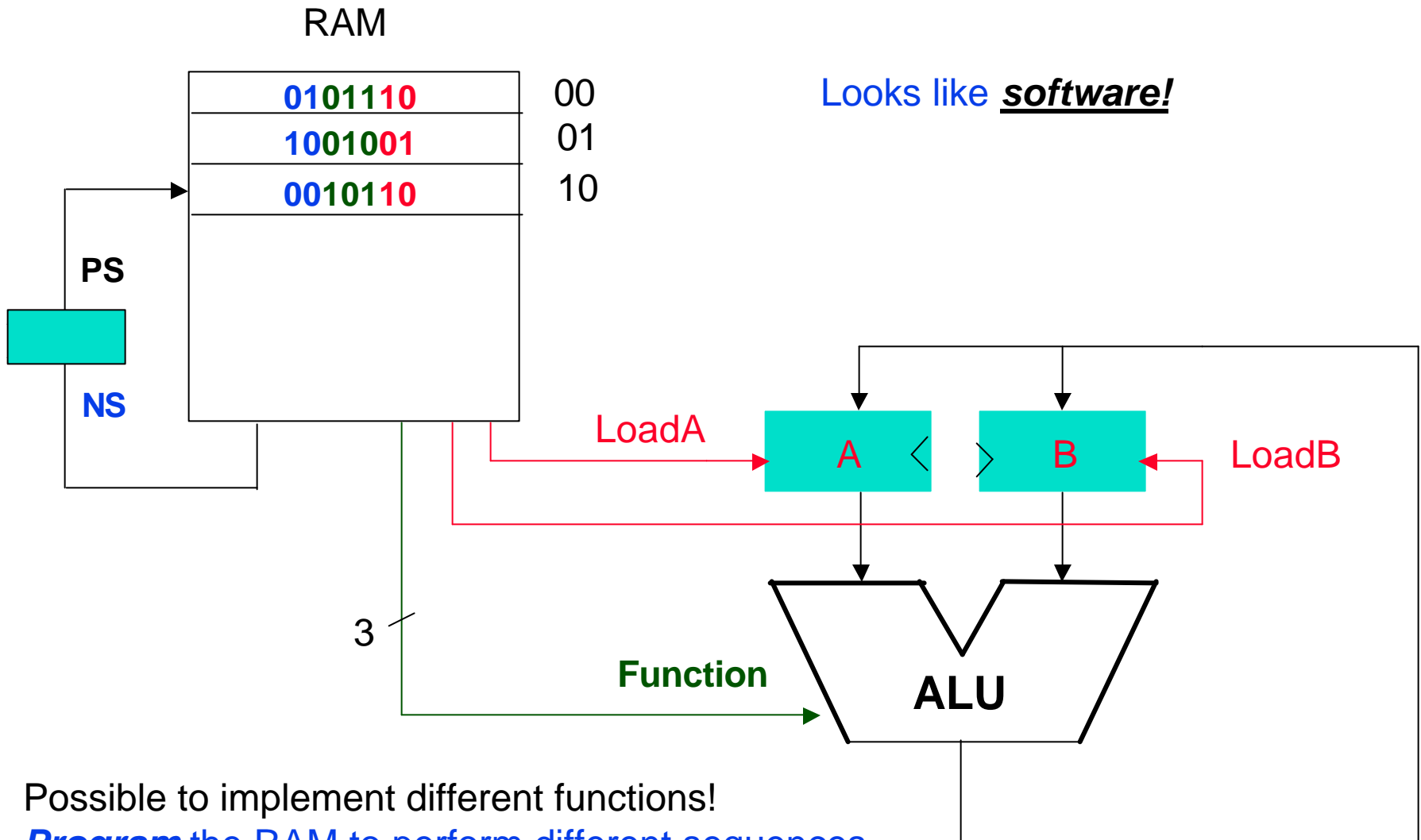


Note: No minimization!
One line in ROM for each state

Putting the Control and Datapath Together



What if we replaced the ROM with RAM?



Summary

- **ALU circuit can perform many functions**
 - **Combinational circuit**
- **ALU chips can be combined together to form larger ALU chips**
 - **Remember to connect carry out to carry in**
- **ALUs form the basis of datapaths**
- **ROMs can form the basis of control paths**
- **Combine the two together to build a computing circuit**
- **Next time: more data and control paths**